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To cite this article: M. Prettì and M. Weigt 2006 *EPL* **75** 8

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Sudden emergence of q -regular subgraphs in random graphs

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received 31 March 2006; accepted 5 May 2006

published online 31 May 2006

PACS. 05.20.-y – Classical statistical mechanics.

PACS. 02.50.-r – Probability theory, stochastic processes, and statistics.

PACS. 89.20.-a – Interdisciplinary applications of physics.

Abstract. – We investigate the computationally hard problem whether a random graph of finite average vertex degree has an extensively large q -regular subgraph, *i.e.*, a subgraph with all vertices having degree equal to q . We reformulate this problem as a constraint-satisfaction problem, and solve it using the cavity method of statistical physics at zero temperature. For $q = 3$, we find that the first large q -regular subgraphs appear discontinuously at an average vertex degree $c_{3\text{-reg}} \simeq 3.3546$ and contain immediately about 24% of all vertices in the graph. This transition is extremely close to (but different from) the well-known 3-core percolation point $c_{3\text{-core}} \simeq 3.3509$. For $q > 3$, the q -regular subgraph percolation threshold is found to coincide with that of the q -core.

Introduction. – In the last years, statistical physics has increasingly been able to analyze and solve complex problems coming from graph theory and theoretical computer science [1, 2]. The interest was particularly focused to so-called *random constraint-satisfaction problems*, which are characterized by a large number of discrete degrees of freedom being subject to an also large number of hard constraints on subsets of variables. The best-known examples are the satisfiability problem, where a set of logical variables is asked to fulfil simultaneously a large number of logical clauses, and the graph-coloring problems, where vertices of a graph are to be assigned colors in a way that no pair of neighboring vertices is equally colored. For both problems, current mathematical tools in discrete mathematics, probability theory, and theoretical computer science do not succeed in solving the models completely. Conversely, new approaches based on the statistical mechanics of disordered systems, in particular the *cavity method* [3], have crucially contributed to our understanding, providing a framework to characterize the statistical properties of the solution space of various constraint-satisfaction problems, and to locate phase transitions in its structure and organization [4, 5].

In this letter, we address a graph-theoretical problem, which at a first glance looks more related to percolation theory than to constraint-satisfaction problems. The question is whether a random graph of given finite connectivity possesses an extensively large q -regular subgraph,

i.e., a subgraph where every vertex has exactly q neighbors (constant degree q). At a closer look, this problem can be naturally embedded into the framework of constraint-satisfaction problems and solved thereby using the cavity method. On the other hand, typical tools from random graph percolation theory are not able to solve the problem. The major reason for this failure is that the problem is NP-complete [6], which means in particular that no linear-time algorithm for searching q -regular subgraphs exists. Mathematical tools based on the analysis of such algorithms, in particular Wormald's rate equation approach [7], are thus unavailable. This is also the major difference with respect to the apparently similar problem of the existence of an extensively large q -core, *i.e.*, the largest subgraph with degrees being equal to *or larger* than q . Such subgraph can be easily found by iteratively removing all vertices of smaller degree. Using the rate equation approach, Pittel, Spencer, and Wormald have shown [8], that such a q -core appears discontinuously at some average random graph degree $c_{q\text{-core}}$, and its size jumps from zero to a finite fraction of all vertices. Let us notice that the existence of an extensive q -core is a necessary condition for a giant q -regular subgraph to exist, since each q -regular subgraph is by definition part of the q -core. Such condition is by no means sufficient, *i.e.*, a q -core may in principle appear before q -regular subgraphs exist at all, so that $c_{q\text{-core}}$ is a lower bound to the emergence of q -regular subgraphs. Bollobás, Kim, and Verstraëte [9], using a refined version of the first-moment method (in statistical physics known better as the *annealed approximation*), have proved that, for $q = 3$, there exists some gap $\gamma > 0$ such that, for $c \in (c_{q\text{-core}}, c_{q\text{-core}} + \gamma)$, *almost surely* no q -regular subgraph exists (*almost surely* means *with probability tending to 1 in the thermodynamic limit* of infinitely large graphs). Moreover, the authors conjecture that the same holds true for any $q > 3$. Looking a bit closer [10] to their proof, it is possible to determine the maximal γ compatible with the first-moment method, which turns out to be as small as $\gamma \simeq 0.0003$ for $q = 3$. This result means that the currently best lower bound for the emergence of 3-regular subgraphs is 3.3512, compared to $c_{3\text{-core}} \simeq 3.3509$ [8]. Conversely, for $q > 3$, the first-moment method cannot improve the lower bound with respect to $c_{q\text{-core}}$, since no positive γ is found.

Are the two transitions really so close to each other? How large is the first q -regular subgraph to appear? How many of these subgraphs are there, and in which way are they related to the q -core? Is the above mentioned conjecture true? Such questions have motivated us to address the problem from the point of view of statistical physics. Using the cavity method, we have found answers, which we argue to be exact.

The model. – Let us start with a more precise definition of the problem. We study random graphs from the Erdős-Rényi ensemble $\mathcal{G}(N, c/N)$ [11]. They have N vertices, and each pair of vertices is connected independently by an edge with probability c/N . The scaling $\mathcal{O}(N^{-1})$ guarantees that the average vertex degree remains finite in the thermodynamic limit $N \rightarrow \infty$, and tends to c . The probability distribution of the degree d approaches a Poissonian of mean c , *i.e.*, $\mathcal{P}_d = e^{-c} c^d / d!$. For this graph ensemble, we ask in general for the existence of q -regular subgraphs. More precisely, we ask whether there exists a *threshold value* $c_{q\text{-reg}}$, such that, in the thermodynamic limit, the probability of finding an extensively large q -regular subgraph tends to 0 for $c < c_{q\text{-reg}}$, and to 1 for $c > c_{q\text{-reg}}$. To answer this, we decorate each edge $\{i, j\}$ with a binary degree of freedom $x_{ij} = x_{ji} \in \{1, 0\}$, meaning, respectively, that the edge is, or is not, in a q -regular subgraph. The constraints are associated to the vertices of the original graph: In each vertex i , either 0 or q “active” links can be present, *i.e.*, $\sum_{j \in \partial i} x_{ij} \in \{0, q\}$, where the sum runs over the set ∂i of all neighbors of vertex i . Note that there is always a solution to these constraints: $x_{ij} \equiv 0 \ \forall \{i, j\}$, *i.e.*, the empty subgraph.

The bipartite structure of variables and constraints can be represented by a suitable factor graph (see fig. 1). Variable nodes are associated to the edges of the original graph, and have

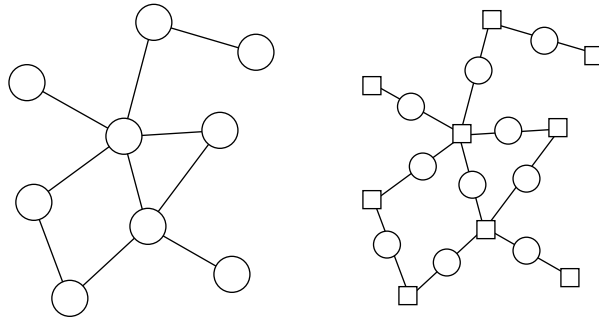


Fig. 1 – Original graph (left) and its factor graph representation (right). The vertices of the original graph become function nodes (squares), whereas the edges of the original graph become decorated by variable nodes (circles).

thus constant degree 2 in the factor graph, whereas function nodes are identified with the original vertices, so that they have the Poissonian degree distribution \mathcal{P}_d of the original graph. We can assign an energy cost to any global configuration $\{x_{ij}\}$ by means of a Hamiltonian counting the number of violated constraints,

$$\mathcal{H}(\{x_{ij}\}) = \sum_{i=1}^N \left[1 - \delta\left(\sum_{j \in \partial i} x_{ij}, 0\right) - \delta\left(\sum_{j \in \partial i} x_{ij}, q\right) \right], \quad (1)$$

with $\delta(\cdot, \cdot)$ denoting a Kronecker delta. Proper q -regular subgraphs are zero-energy ground states of this Hamiltonian, and their properties can be analyzed using the cavity method at zero temperature. The main task is to calculate the quenched entropy density

$$s = \lim_{N \rightarrow \infty} N^{-1} \overline{\ln Z}, \quad (2)$$

where the zero-temperature partition function,

$$Z = \sum_{\{x_{ij}\}} \delta(\mathcal{H}, 0), \quad (3)$$

represents the number of q -regular subgraphs, while the overline denotes the average over the random graph ensemble. The phase transition point $c_{q\text{-reg}}$ can be identified as the average vertex degree where this entropy first takes a positive value.

Note that Hamiltonian (1) contains two special cases which were recently addressed with very similar methods. First, for $q = 1$, the problem of matchings in random graphs is recovered, cf. [12, 13]. Second, for $q = 2$, the regular subgraphs are loops [14–16]. There is, however, one big difference: For $q \leq 2$, the existence problem becomes polynomially solvable, and the resulting physical picture of the phase transition is simpler. In a random graph, extensive matchings exist whenever there is an extensive number of links ($c > 0$), and extensive loops appear continuously at the random-graph percolation transition [11] at $c = 1$. The hard problem addressed in the afore-mentioned papers is therefore the *counting problem*, whereas, for $q \geq 3$, even asking about the very existence of q -regular subgraphs is NP-complete.

The replica-symmetric solution. – The central step of the cavity method is to set up a *message passing* procedure, incorporating the local consequences of the degree constraints

into a globally self-consistent iterative scheme. We denote by $p_{i \rightarrow j}$ an elementary message, *i.e.*, the probability that constraint i forces edge $\{i, j\}$ to be present in the subgraph. This happens if and only if exactly $q - 1$ other edges incident to i are present. We can therefore write

$$p_{i \rightarrow j} = \frac{w_{\partial i \setminus j \rightarrow i}^{q-1}}{\sum_{n=0, q-1, q} w_{\partial i \setminus j \rightarrow i}^n}, \quad (4)$$

where we have defined

$$w_{V \rightarrow i}^n \equiv \sum_{\substack{U \subseteq V \\ |U|=n}} \prod_{j \in U} p_{j \rightarrow i} \prod_{k \in V \setminus U} (1 - p_{k \rightarrow i}). \quad (5)$$

According to (5), the numerator of eq. (4) counts the joint probability that exactly $q - 1$ of the edges arriving in i from other vertices than j (*i.e.*, from vertices in the set $\partial i \setminus j$) are forced to be in the subgraph by their second end-vertex. The denominator sums over all possibilities to have 0, q , or $q - 1$ such incoming edges. Note that only these three possibilities are consistent with the constraint, *i.e.*, normalization explicitly excludes contradictory situations. Note also that the joint probability is assumed to factorize in the edges, which, in the replica-symmetric situation (one thermodynamic state only [3]), is expected to become exact for $N \gg 1$, due to the locally tree-like organization of random graphs. As a last remark, we note that $p_{i \rightarrow j} \equiv 0$ is always a solution of eqs. (4,5), corresponding to the empty subgraph, which trivially satisfies all degree constraints.

Having found a fixed point of these equations, we calculate observables like the probabilities p_i , p_{ij} that a vertex or an edge, respectively, are in a q -regular subgraph,

$$p_i = \frac{w_{\partial i \rightarrow i}^q}{\sum_{n=0, q} w_{\partial i \rightarrow i}^n}, \quad (6)$$

$$p_{ij} = \frac{p_{i \rightarrow j} p_{j \rightarrow i}}{p_{i \rightarrow j} p_{j \rightarrow i} + (1 - p_{i \rightarrow j})(1 - p_{j \rightarrow i})}. \quad (7)$$

The numerator of eq. (6) counts the probability that q edges arriving in vertex i are forced to be in the subgraph, and it has to be normalized by the sum over all consistent possibilities: Either q edges (vertex in the subgraph) or 0 edges (vertex not in the subgraph) can be in the subgraph. The numerator of eq. (7) contains the case that the messages coming from the end-vertices of edge $\{i, j\}$ consistently force the edge to be element of the subgraph, and it has to be normalized with respect to all consistent message pairs. Further on, we can calculate the entropy, which results from eq. (2) via

$$\ln Z = \sum_{i=1}^N \ln z_i - \sum_{\{i, j\}} \ln z_{ij}, \quad (8)$$

where z_i and z_{ij} denote respectively the denominators of the right hand sides of eqs. (6) and (7).

Equations (4), (5) can be solved either directly on single graph instances, or in distribution in the average over the random graphs. Due to the random structure of the underlying graph, these equations are easily translated into a self-consistent equation for the message distribution $\rho(p)$,

$$\rho(p) = \sum_{d=0}^{\infty} e^{-c} \frac{c^d}{d!} \int_0^1 dp_1 \rho(p_1) \dots \int_0^1 dp_d \rho(p_d) \delta(p - f_d(p_1, \dots, p_d)), \quad (9)$$

where $\delta(\cdot)$ denotes a Dirac delta, and with f_d being given by eqs. (4), (5). In complete analogy, we can also derive equations for the ensemble-averaged distribution of true occupation probabilities p_i, p_{ij} and for the entropy. The trivial solution $\rho(p) = \delta(p)$ exists independently of c , and has zero entropy. The question of the existence of extensive q -regular subgraphs reduces to the question of the existence and thermodynamic stability of non-trivial solutions $\rho(p)$. A full solution can be constructed only numerically, using a population-dynamical scheme [3], but important information about the onset of a non-trivial solution, and its relation to the q -core, can be read off analytically from eq. (9). To do so, we first simplify it by projecting the real-valued probabilities p to a ternary variable $X = 0, 1, *$, depending on whether $p = 0, 1$, or $0 < p < 1$. The corresponding weights in $\rho(p)$ are denoted by P_X . The trivial solution has of course $P_0 \equiv 1$ and $P_1 \equiv P_* \equiv 0$, whereas in general we have to derive closed equations for the three probabilities P_X , for $X = 0, 1, *$. This task becomes considerably simplified by the observation that P_1 must vanish, since, in the thermodynamic limit, a finite fraction of messages polarized to $p = 1$ would necessarily lead to contradictions to the degree constraints. Moreover, it turns out that a nontrivial ($X = *$) out-message is sent if and only if $q - 1$ or more in-messages are also nontrivial. This translates to

$$P_* = \sum_{d=q-1}^{\infty} e^{-c} \frac{c^d}{d!} \sum_{n=q-1}^d \binom{d}{n} P_*^n P_0^{d-n} = e^{-cP_*} \sum_{n=q-1}^{\infty} \frac{(cP_*)^n}{n!}, \quad (10)$$

where we have eliminated P_0 using normalization $P_0 + P_* = 1$. This equation is not new: It appears in the context of the q -core, and its first non-trivial solution $P_* > 0$ appears exactly at the q -core threshold [8]. Here, we find this connectivity as the *spinodal point* for q -regular subgraphs: The first non-trivial solution of eq. (9) appears at this point, but its thermodynamic stability is not guaranteed in principle. This observation also clarifies the relation between the q -core and all possible q -regular subgraphs: Edges not belonging to the q -core never belong to any q -regular subgraph, whereas almost all edges in the q -core belong to some but not all q -regular subgraphs.

Unfortunately, the projection to ternary variables does not allow for the calculation of the entropy, which is the important thermodynamic potential for checking which solution of eq. (9) is actually the thermodynamically stable one. As previously mentioned, the entropy depends on the full information carried by $\rho(p)$, so that we have to analyze eq. (9) numerically. We have done so, using a representation of $\rho(p)$ via a population of 2^{20} elements, and have used a usual iterative update scheme [3].

In fig. 2, we report the results in terms of average subgraph size (fraction ϕ of vertices in the subgraph) and quenched entropy density s , for the case $q = 3$. A non-trivial solution appears at $c_{3\text{-core}}$, but it has negative entropy. The thermodynamically stable solution remains the trivial, zero-entropy one, and no extensive 3-regular subgraph exists. Nevertheless, we can see that the entropy increases upon growing average degree c , and becomes positive at $c_{3\text{-reg}} \simeq 3.3546$, where a *first-order phase transition* takes place. The first 3-regular subgraphs appear discontinuously, containing immediately about 24% of the full graph. The situation is different for larger q values. In particular, we have investigated the cases $q = 4, 5$, for which one can observe that the non-trivial solution still appears at $c_{q\text{-core}}$, but with an already positive entropy (see fig. 3). We are thus led to conclude that, contrary to the conjecture of Bollobás and coworkers [9], the emergence of extensive q -regular subgraphs coincides with that of the q -core, for $q > 3$, whereas $q = 3$ is a peculiar case.

Stability of replica symmetry. – So far, our results rely on the assumption of replica symmetry. Does the inclusion of possible replica-symmetry-broken solutions change the threshold

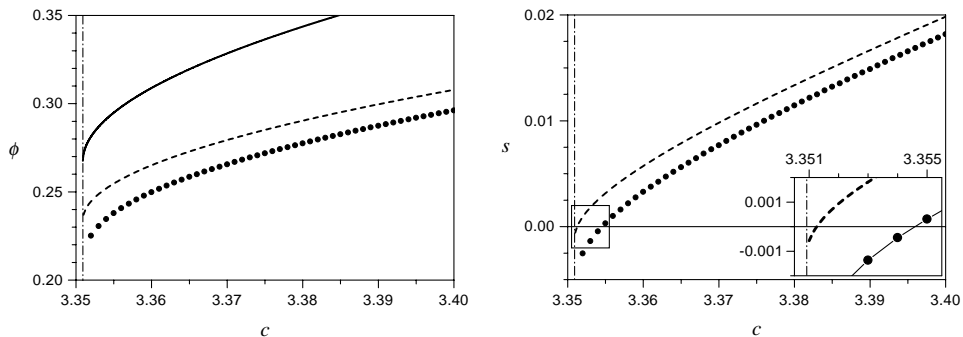


Fig. 2 – Average subgraph size ϕ (left) and entropy density s (right) as a function of the average vertex degree c , for 3-regular subgraphs (circles) and for the 3-core (solid lines). Dashed lines denote results from the annealed approximation for 3-regular subgraphs [10], which provides an upper bound for the entropy, and a lower bound for the threshold. Thin dash-dotted lines mark the 3-core threshold.

value? In order to investigate this issue, we have set up the cavity equations for one-step replica symmetry breaking (1-RSB) [3]; details will be given elsewhere [10]. Using these equations, we have first verified the local stability of the replica-symmetric solution described above, *i.e.*, that small perturbations decay exponentially fast. Moreover, we have searched without success for a non-trivial, locally stable 1-RSB solution, which could possibly appear discontinuously. Based on these findings, we conjecture that the replica-symmetric values $c_{3\text{-reg}} \simeq 3.3546$, and $c_{q\text{-reg}} = c_{q\text{-core}}$ for $q > 3$, are exact. The mathematical proof of these statements remains an interesting open problem. Let us notice that, given the replica-symmetric character of the transition, such a proof might be more easily accessible than the proofs of exactness of statistical-physics results for threshold phenomena in other constraint-satisfaction problems.

Conclusion and outlook. – In this letter, we have analyzed the emergence (percolation) of q -regular subgraphs in random graphs. Using the cavity method of statistical physics, we have found that this happens in a first-order transition, *i.e.*, the subgraphs are immediately extensively large. These results are based on a replica-symmetric calculation. Nevertheless, stability with respect to replica symmetry breaking leads us to conjecture that the observed threshold values are exact. For $q = 3$, the transition occurs in extreme vicinity to (but deviates from) the well-known 3-core percolation point $c_{3\text{-core}} \simeq 3.3509$, whereas, for larger q values

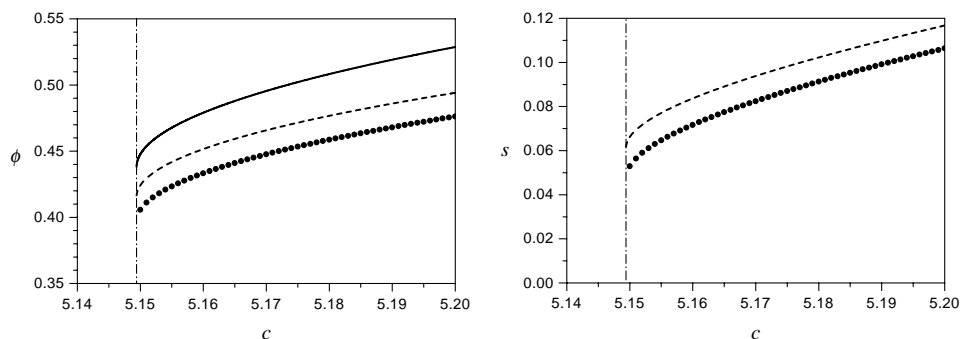


Fig. 3 – The same as fig. 2 for $q = 4$.

the threshold is found to coincide with that of the q -core. Moreover, our method clarifies the relationship between these apparently similar, but computationally very different problems: Whenever q -regular subgraphs exist, their union equals the q -core.

In a future publication [10], we shall further elucidate the structure of such q -regular subgraphs. Our method, as formulated here, allows to identify the entropy and thus also the size of the *most frequent* q -regular subgraphs. One can go beyond this by coupling the subgraph size to a conjugate chemical potential, and analyze smallest and largest q -regular subgraphs. Whereas it is mathematically clear that the q -core, as the maximal subgraph of minimal degree at least q , is unique, a similar statement does not exist for maximal q -regular subgraphs. A further interesting question is the one for conditions of existence of a q -factor (*i.e.* a spanning q -regular subgraph) in random graphs of given degree sequence. According to a conjecture of Bollobás *et al.*, these are expected to exist if the minimal degree in the original graph is $q + 1$.

* * *

We are indebted to G. SEMERJIAN for many interesting discussions, and for communicating to us his results prior to publication. We also acknowledge helpful discussions with M. LEONE, A. PAGNANI, and L. ZDEBOROVÁ.

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